## Outline

- Chapter 2: Direct Link Networks
- Encoding
- Framing
- Error Detection
- Sliding Window Algorithm


## Direct Link Networks

## Direct Link Networks

- Hosts are directly connected by some medium
- Twisted pair: telephone cable, Ethernet (Category 5: Cat5)
- Coaxial pair: TV
- Optical Fiber
- Wireless: Infrared, Radio, Microwave
- Common bandwidth designators:
- DS1 (or T1): 1.544 Mbps
- DS3 (or T3): 44.736 Mbps (for example, Charter Athens has 2 DS3 links now)
- STS-1 (OC1): 51.840 Mbps
- STS-12: 622.080 Mbps ...


## Last Mile

- Plain Old Telephone Service) POTS:
- 28.8 Kbps to 56 Kbps
- ISDN
- xDSL 1.544 Mbps to 8.448 Mbps
- Cable (40 Mbps down, 20 Mbps up) - Shared
- wish we can get that much huh?


## Encoding

## Encoding

- Signals propagate over a physical medium
- modulate electromagnetic waves
- e.g., vary voltage
- Encode binary data onto signals
- e.g., 0 as low signal and 1 as high signal
- known as Non-Return to zero (NRZ)



## Problem: Consecutive 1s or 0s

- Low signal (0) may be interpreted as no signal - High signal (1) leads to baseline wander - Unable to recover clock


## Alternative Encodings

- Non-return to Zero Inverted (NRZI)
- make a transition from current signal to encode a one; stay at current signal to encode a zero
- solves the problem of consecutive ones
- Manchester
- transmit XOR of the NRZ encoded data and the clock
- only 50\% efficient.


## Encodings (cont)

- 4B/5B
- every 4 bits of data encoded in a 5-bit code
- 5 -bit codes selected to have no more than one leading 0 and no more than two trailing 0s
- thus, never get more than three consecutive 0s
- resulting 5-bit codes are transmitted using NRZI
- achieves 80\% efficiency


## Encodings (cont)



## Framing

## Framing

- Break sequence of bits into a frame
- Typically implemented by network adaptor


Frames

## Approaches

- Sentinel-based
- delineate frame with special pattern: 01111110
- e.g., HDLC (ISO), SDLC (IBM), PPP (dialup)

| 16 |  |  |  |
| :--- | :---: | :--- | :--- | :--- |
| Beginning <br> sequence | Header | Body |  |

- problem: what if the special pattern appears in the payload itself?
- solution: bit stuffing
- sender: insert 0 after five consecutive 1s
- receiver: delete 0 that follows five consecutive 1s


## Approaches (cont)

- Counter-based
- include payload length in header
- e.g., DDCMP (DECNET)

- problem: count field itself corrupted
- solution: catch when CRC fails


## Approaches (cont)

- Clock-based
- each frame is 125us long
- e.g., SONET: Synchronous Optical Network
- STS-n (STS-1 = 51.84 Mbps)

Overhead


## Error Detection

## Cyclic Redundancy Check

- Add k bits of redundant data to an n -bit message
- want $k \ll n$
- e.g., $k=32$ and $n=12,000$ (1500 bytes)
- Represent n -bit message as n - 1 degree polynomial
- e.g., MSG=10011010 as $M(x)=x 7+x 4+x 3+x 1$
- Let k be the degree of some divisor polynomial
- e.g., $C(x)=x 3+x 2+1$


## CRC (cont)

- Transmit polynomial $\mathrm{P}(\mathrm{x})$ that is evenly divisible by C(x)
- shift left k bits, i.e., $M(x) x k$
- subtract remainder of $M(x) x k / C(x)$ from $M(x) x k$
- Receiver polynomial $\mathrm{P}(\mathrm{x})+\mathrm{E}(\mathrm{x})$
$-E(x)=0$ implies no errors
- Divide $(P(x)+E(x))$ by $C(x)$; remainder zero if:
- $E(x)$ was zero (no error), or
$-E(x)$ is exactly divisible by $C(x)$


## Selecting C(x)

- All single-bit errors, as long as the xk and x0 terms have non-zero coefficients.
- All double-bit errors, as long as $\mathrm{C}(\mathrm{x})$ contains a factor with at least three terms
- Any odd number of errors, as long as $C(x)$ contains the factor $(x+1)$
- Any 'burst' error (i.e., sequence of consecutive error bits) for which the length of the burst is less than $k$ bits.
- Most burst errors of larger than k bits can also be detected
- See Table 2.6 on page 102 for common $C(x)$


## Internet Checksum Algorithm

- View message as a sequence of 16 -bit integers; sum using 16-bit ones-complement arithmetic; take ones-complement of the result.
u_short cksum(u_short *buf, int count) \{ register u_long sum $=0$; while (count--) \{ sum += *buf++; if (sum \& OxFFFFO000) \{
/* carry occurred, so wrap around */ sum \&= OxFFFF; sum++;
\}
\}
return ~(sum \& OxFFFF);
\}

