Overview

Store-and-Forward Switches

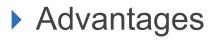
- Direct link network technologies do not scale. You cannot add all the nodes in the world into a single direct link network.
 - If all the nodes in the world were to broadcast, then there is chaos.
 - If all the nodes in the world are in a switched network, then each node will get infinitesimally small quanta
 - If all the nodes in the world were using point-to-point, then we need to run a whole bunch of wires between each node
- Solution: Partition the network into multiple zones and connect each one using a switch. The switch will forward packets from input to output. Only packets that need to be forwarded are sent; local packets are not forwarded

Scalable Networks

Switch

- forwards packets from input port to output port
 - Each input/output can use different "direct link network"
- port selected based on address in packet header

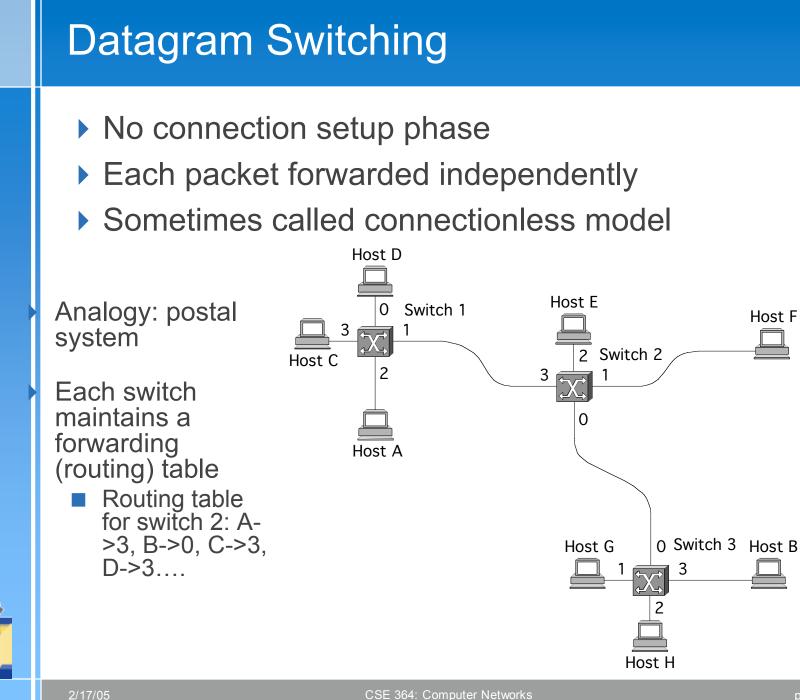




- cover large geographic area (tolerate latency)
- support large numbers of hosts (scalable bandwidth)

Design goals

- Switches should increase aggregate throughput
- Switches look inside the packet to decide which output port it should place the packet in
 - Switches look at the destination address
- Switches number input and output ports so that they know where to send the packet
- Switching strategies:
 - Datagram: packet carries enough information, along with forwarding/routing tables
 - Virtual circuit: connection establishment phase
 - Source routing: packet carries the routing table

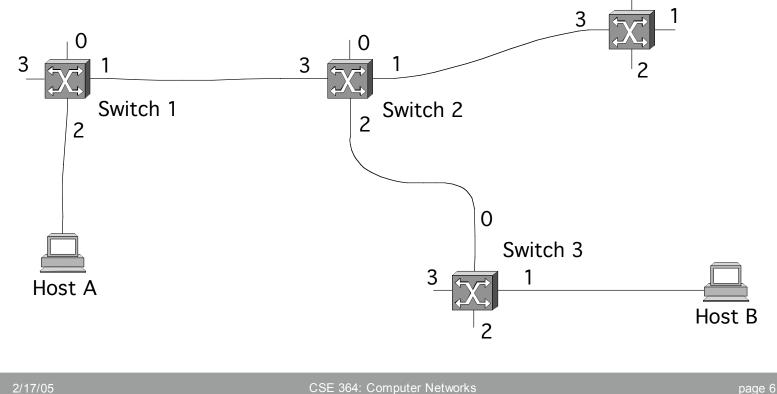


Datagram Model

- There is no round trip time delay waiting for connection setup; a host can send data as soon as it is ready
- Source host has no way of knowing if the network is capable of delivering a packet or if the destination host is even up.
- Since packets are treated independently, it is possible to route around link and node failures
- Since every packet must carry the full address of the destination, the overhead per packet is higher than for the connection-oriented model

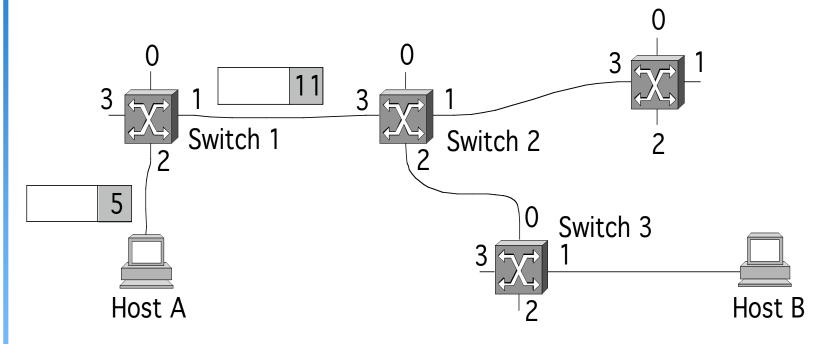
Virtual Circuit Switching

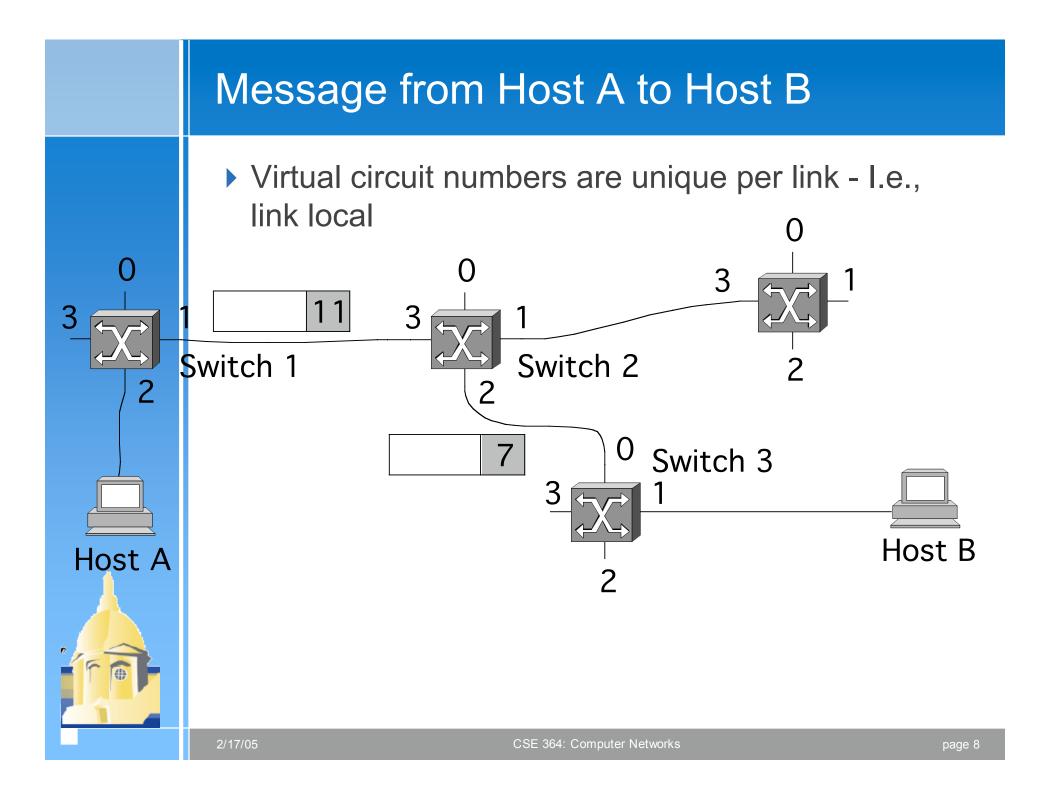
- Explicit connection setup (and tear-down) phase
- Subsequence packets follow same circuit
- Sometimes called connection-oriented model
- Analogy: phone call



Message from Host A to Host B

- Each switch maintains a VC table
 - Incoming port 2, Incoming VC 5, Outgoing Interface 1, Outgoing VC 11





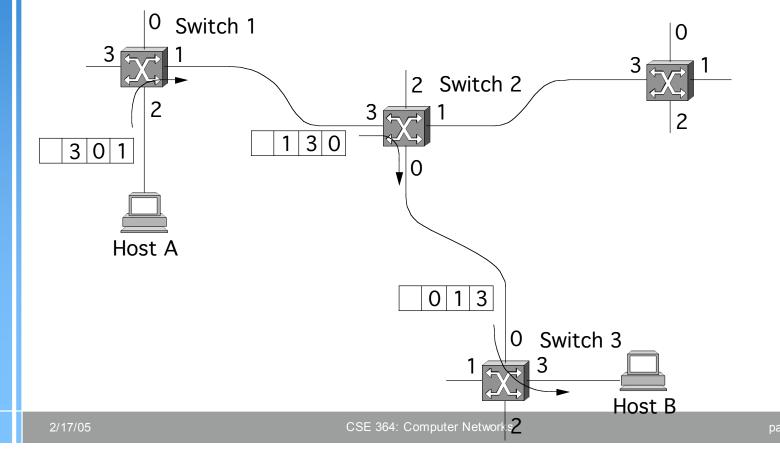
Virtual Circuit Model

- Typically wait full RTT for connection setup before sending first data packet.
- While the connection request contains the full address for destination, each data packet contains only a small identifier, making the per-packet header overhead small.
- If a switch or a link in a connection fails, the connection is broken and a new one needs to be established.
- Connection setup provides an opportunity to reserve resources.

Source Routing

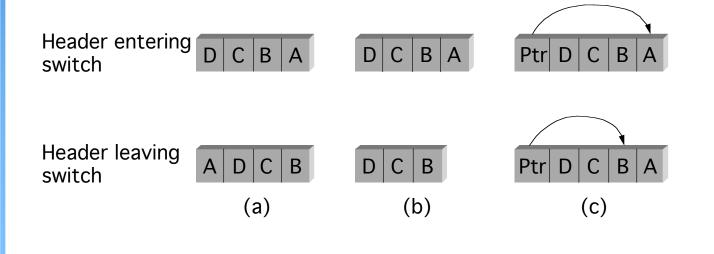
Each packet carries the routing information

- Source host knows the exact route
- Doesn't scale, but reduces state in switch
- Node rotate the address



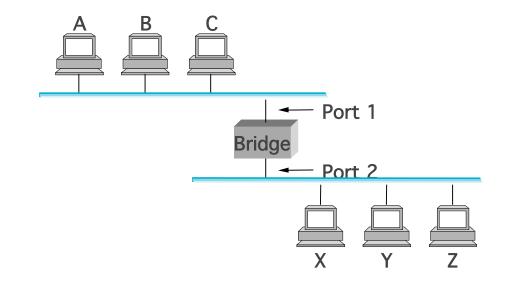
Maintaining the source route

- A. rotation of route
 - We will know the reverse path
- B. Stripping
 - Packets become smaller
- C. Pointer
 - Each router updates the pointer



Bridges and Extended LANs

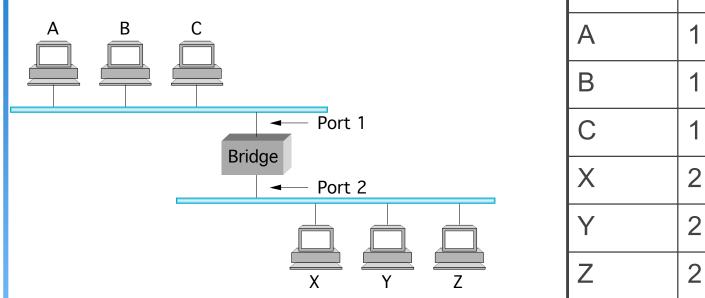
- LANs have physical limitations (e.g., 2500m)
- Connect two or more LANs with a bridge
 - accept and forward strategy
 - level 2 connection (does not add packet header)



Ethernet Switch = Bridge on Steroids

Learning Bridges

- Do not forward when unnecessary
- Maintain forwarding table



- Learn table entries based on source address
- Table is an optimization; need not be complete
- Always forward broadcast frames

Port

Host

Lab setup

Gateway13-18 are connected to a 8-port Gigabit switch. Itanium servers are connected to a 5 port gigabit switch. The uplink from the 5 port switch is connected to the 8-port switch. The uplink from the 8-port switch is connected to another switch which is connected to ND

