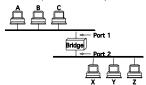
Bridges and Extended LANs

- LANs have physical limitations (e.g., 2500m)
- · Connect two or more LANs with a bridge
 - accept and forward strategy
 - level 2 connection (does not add packet header)



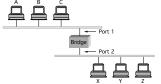
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• Ethernet Switch = Bridge on Steroids



Learning Bridges

- · Do not forward when unnecessary
- · Maintain forwarding table



Host	Port
Α	1
В	1
С	1
Χ	2
Υ	2
Z	2

- Learn table entries based on source address
- · Table is an optimization; need not be complete
- · Always forward broadcast frames



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Spanning Tree Algorithm

Problem: loops - no mechanism to remove looping frames







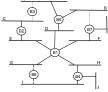
- · Bridges run a distributed spanning tree algorithm
- select which bridges actively forward
 - developed by Radia Perlman
 - now IEEE 802.1 specification

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Algorithm Overview

- Each bridge has unique id (e.g., B1, B2, B3)
- · Select bridge with smallest id as root
- · Select bridge on each LAN closest to root as designated bridge (use id to break ties)
 - Each bridge forwards frames over each LAN for which it is the designated bridge



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Algorithm Details

- · Bridges exchange configuration messages
 - id for bridge sending the message
 - id for what the sending bridge believes to be root bridge
 - distance (hops) from sending bridge to root bridge
- · Each bridge records current best configuration message for each port
- · Initially, each bridge believes it is the root

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Algorithm Detail (cont)

- · When learn not root, stop generating config messages
 - in steady state, only root generates configuration messages
- · When learn not designated bridge, stop forwarding config messages
 - in steady state, only designated bridges forward config messages
- · Root continues to periodically send config messages
- · If any bridge does not receive config message after a period of time, it starts generating config messages claiming to be the root



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Broadcast and Multicast · Forward all broadcast/multicast frames - current practice · Learn when no group members downstream · Accomplished by having each member of group G send a frame to bridge multicast address with G in source field Feb-12-04 4/598N: Computer Networks **Tcpdump trace** · tcpdump -p 02:21:52.651816 802.1d config 0000.00:02:2d:71:03:ef.0001 root 0000.00:02:2d:71:03:ef pathcost 0 age 0 max 20 helio 2 fdelay 15 02:21:53.263956 engr-fe21.gw.nd.edu > ALL-SYSTEMS.MCAST.NET: igmp query v2 [tos 0xc0] [til 1] 02:25:22.656898 CDP v2, ttl=180s DevID '013183892(hub24-1b.hub.nd.edu)' Addr (1): IPv4 129.74.24.67 PortID '5/10' CAP 0x0e[|cdp] Feb-12-04 4/598N: Computer Networks **Limitations of Bridges** · Do not scale - spanning tree algorithm does not scale - traffic gets bridged through the root bridge · Spanning tree is designed to avoid loops, not traffic balancing: redundant routes are ignored - broadcast does not scale · Do not accommodate heterogeneity · Caution: beware of transparency Feb-12-04

VLAN Create virtual lans (broadcast domains) without rewiring Add a 4 byte VLAN id to each frame between P routing and bridging TRANS SmartBridges http://www.uwtv.org/programs/displayevent.asp?rid=782 Hybrid between IP routing and bridging

Feb-12-04